

Throughout this note we fix a symmetric monoidal category $(\mathcal{C}, \otimes, \tau)$ where τ is the symmetry functor with $\tau_{AB} : A \otimes B \rightarrow B \otimes A$.

A bicommutative bialgebra object in \mathcal{C} is an object A , with functors $\mu_A : A \otimes A \rightarrow A$, $\eta_A : I \rightarrow A$, $\Delta_A : A \rightarrow A \otimes A$, and $\epsilon_A : A \rightarrow I$ called multiplication, unit, comultiplication and counit respectively, satisfying certain axioms.

Given two bicommutative bialgebra objects A and B in \mathcal{C} , we may give $A \otimes B$ the structure of a bicommutative bialgebra: $\mu_{A \otimes B} := (\mu_A \otimes \mu_B)(1 \otimes \tau_{BA} \otimes 1)$, $\eta_{A \otimes B} = \eta_A \otimes \eta_B$, $\Delta_{A \otimes B} = (1 \otimes \tau_{AB} \otimes 1)(\Delta_A \otimes \Delta_B)$, and $\epsilon_{A \otimes B} = \epsilon_A \otimes \epsilon_B$. In this way we may make $A^{\otimes m}$ into a bicommutative bialgebra for each $m \geq 2$.

As a point of notation, we will write μ_A^n for the functor from $A^{\otimes n} \rightarrow A$, given inductively by $\mu_A^0 = \eta_A$, and $\mu_A^{n+1} = \mu(1 \otimes \mu_A^n)$, and Δ_A^n for the functor from $A \rightarrow A^{\otimes n}$ given inductively by $\Delta_A^0 = \epsilon_A$ and $\Delta_A^{n+1} = \Delta_A(1 \otimes \Delta_A^n)$. Notice every μ_A^n and Δ_A^n is bialgebra map.

Lemma. *The monoidal subcategory of \mathcal{C} whose objects are bicommutative bialgebra objects in \mathcal{C} and whose morphisms are bialgebra bicomorphisms may be enriched over commutative monoids.*

From now on we will refer to this enriched category as \mathbf{Bialg} . Notice that the enriched structure makes $\text{Hom}_{\mathbf{Bialg}}(A, A)$ is a rig with identity id_A whenever A is an object in \mathbf{Bialg} . We will just write $\text{End}(A)$ for this rig.

Proof. We first need to explain how to define an addition on the Hom sets. Suppose A and B are bicommutative bialgebras, and suppose that f and g are two morphisms from A to B . We define

$$f + g := \mu_B(f \otimes g)\Delta_A$$

For each Hom set the axioms for a commutative monoid now follow easily, but for completeness: suppose f, g and h are in $\text{Hom}(A, B)$

$$\begin{aligned} (f + g) + h &= \mu_B((\mu_B(f \otimes g)\Delta_A) \otimes h)\Delta_A \\ &= \mu_B(\mu_B \otimes 1)(f \otimes g \otimes h)\Delta_A(1 \otimes \Delta_A) \\ &= \mu_B(1 \otimes \mu_B)(f \otimes g \otimes h)(1 \otimes \Delta_A)\Delta_A \\ &= f + (g + h) \end{aligned}$$

where the third equality follows from the coassociativity of A , and associativity of B .

The zero map from A to B is the composite of the counit ϵ_A and the unit η_B , and

$$0 + f = \mu_B(\eta_B \epsilon_A \otimes f)\Delta_A = \mu_B(\eta_B \otimes 1)(1 \otimes f)(\epsilon_A \otimes 1)\Delta_A = f$$

where the last equality follows from the axioms for unit and counit.

The symmetry of $+$ follows from the commutativity of B , and cocommutativity of A : if τ_A is the symmetry map $A \otimes A \rightarrow A \otimes A$ then

$$f + g = \mu_B(f \otimes g)\Delta_A = \mu_B \tau_B(f \otimes g)\Delta_A = \mu_B(g \otimes f)\tau_A \Delta_A = \mu_B(g \otimes f)\Delta_A = g + f$$

To complete the proof we need to check that the composition of morphisms gives a monoid map $\text{Hom}(B, C) \times \text{Hom}(A, B) \rightarrow \text{Hom}(A, C)$: suppose f_1 and f_2 are in $\text{Hom}(B, C)$ and g_1 and g_2 are in $\text{Hom}(A, B)$ then

$$f_1(g_1 + g_2) = f_1 \mu_B(g_1 \otimes g_2)\Delta_A = \mu_C(f_1 \otimes f_1)(g_1 \otimes g_2)\Delta_A = \mu_C(f_1 g_1 \otimes f_1 g_2)\Delta_A = f_1 g_1 + f_1 g_2$$

and

$$(f_1 + f_2)g_1 = \mu_C(f_1 \otimes f_2)\Delta_B g_1 = \mu_C(f_1 \otimes f_2)(g_1 \otimes g_1)\Delta_A = \mu_C(f_1 g_1 \otimes f_2 g_1)\Delta_A = f_1 g_1 + f_2 g_1.$$

□

Our goal now is the following theorem:

Theorem. *If R is a rig then $\text{Mat}(R)$ is just the PROP for bicommutative bialgebras A equipped a map of rigs $R \rightarrow \text{End}(A)$.*

We prove this theorem with three lemmas. Firstly we show

Lemma. *If A is a bicommutative bialgebra and $\phi : R \rightarrow \text{End}(A)$ a map of rigs, then there is a strict monoidal functor F_A enriched over commutative monoids from $\text{Mat}(R)$ to Bialg such that $F_A(1) = A$ and F_A is just ϕ on $\text{Hom}(1, 1)$.*

Proof. Because F_A is a strict monoidal functor with $F_A(1) = A$, $F_A(n)$ is necessarily $A^{\otimes n}$ for every n .

We begin by defining F_A on $\text{Hom}(n, 1)$ for each n . If (r_i) is an $(1 \times n)$ matrix with entries in R we set $F_A(r_i) = \mu_A^n(\phi(r_1) \otimes \cdots \otimes \phi(r_n))$. Notice that in particular $F_A(r : 1 \rightarrow 1)$ is just $\phi(r)$ as required. Also notice $F_A(0 : 0 \rightarrow 1) = \eta_A$ and $F_A(11 : 2 \rightarrow 1)$ is just μ_A . We need all these maps $F_A(r_i)$ to be bialgebra maps, but this is true because they are defined as a composite of bialgebra maps.

Now suppose that (r_{ij}) is any $(m \times n)$ matrix with entries in R . We define $F_A(r_{ij}) = (F_A(r_{i1}) \otimes \cdots \otimes F_A(r_{in}))\Delta_{A^{\otimes n}}^m$. This time it may be easily seen that $F_A(0 : 1 \rightarrow 0) = \epsilon_A$ and $F_A((11)^t : 1 \rightarrow 2) = \Delta_A$. As before all these maps $F_A(r_{ij})$ are bialgebra maps because they are a composite of such.

We need to check that F as defined is an enriched functor. First, we check that $F_A(r_{ij} + s_{ij}) = F_A(r_{ij}) + F_A(s_{ij})$ for every pair of R -valued $(m \times n)$ matrices (r_{ij}) and (s_{ij}) . As before we begin by considering the case $m = 1$ suppressing the second index as we may:

$$\begin{aligned} F_A(r_i + s_i) &= \mu_A^n(\phi(r_1 + s_1) \otimes \cdots \otimes \phi(r_n + s_n)) \\ &= \mu_A^n((\phi(r_1) + \phi(s_1)) \otimes \cdots \otimes (\phi(r_n) + \phi(s_n))) \\ &= \mu_A^n((\mu_A(\phi(r_1) \otimes \phi(s_1))\Delta_A) \otimes \cdots \otimes (\mu_A(\phi(r_n) \otimes \phi(s_n))\Delta_A)) \\ &= \mu_A^n(\mu_A \otimes \cdots \otimes \mu_A)(\phi(r_1) \otimes \phi(s_1) \otimes \cdots \otimes \phi(r_n) \otimes \phi(s_n))(\Delta_A \otimes \cdots \otimes \Delta_A) \\ &= \mu_A^n(\mu_{A^{\otimes n}}(\phi(r_1) \otimes \cdots \otimes \phi(r_n) \otimes \phi(s_1) \otimes \cdots \otimes \phi(s_n))\Delta_{A^{\otimes n}}) \\ &= \mu_A(\mu_A^n(\phi(r_1) \otimes \cdots \otimes \phi(r_n)) \otimes \mu_A^n(\phi(s_1) \otimes \cdots \otimes \phi(s_n)))\Delta_{A^{\otimes n}} \\ &= F_A(r_i) + F_A(s_i) \end{aligned}$$

Now we consider the general case:

$$\begin{aligned} F_A(r_{ij} + s_{ij}) &= (F_A(r_{i1} + s_{i1}) \otimes \cdots \otimes F_A(r_{in} + s_{in}))\Delta_{A^{\otimes n}}^m \\ &= ((F_A(r_{i1}) + F_A(s_{i1})) \otimes \cdots \otimes (F_A(r_{in}) + F_A(s_{in})))\Delta_{A^{\otimes n}}^m \\ &= (\mu_A(F_A(r_{i1}) \otimes F_A(s_{i1}))\Delta_{A^{\otimes n}}) \otimes \cdots \otimes (\mu_A(F_A(r_{in}) \otimes F_A(s_{in}))\Delta_{A^{\otimes n}}))\Delta_{A^{\otimes n}}^m \\ &= \mu_{A^{\otimes m}}((F_A(r_{i1}) \otimes \cdots \otimes F_A(r_{in}))\Delta_{A^{\otimes n}}^m) \otimes (F_A(s_{i1}) \otimes \cdots \otimes F_A(s_{in}))\Delta_{A^{\otimes n}}^m)\Delta_{A^{\otimes n}} \\ &= \mu_{A^{\otimes m}}(F_A(r_{ij}) \otimes F_A(s_{ij}))\Delta_{A^{\otimes n}} \\ &= F_A(r_{ij}) + F_A(s_{ij}) \end{aligned}$$

We can now complete the proof by showing that F_A is a functor i.e. that F_A preserves composition. Because we have already checked that F_A preserves the

enriched structure it suffices to check this on matrices with precisely one non-zero entry since these generate all matrices under $+$.

But if we have a $(1 \times n)$ row matrix with all entries zero except possibly the i th entry which takes value r , it is easy to check that F_A sends it to the map $\epsilon_{A^{\otimes i-1}} \otimes \phi(r) \otimes \epsilon_{A^{\otimes n-i}}$. Then we see that F_A sends a general $(m \times n)$ matrix with all entries 0 except the ij th which takes value r to $\eta_{A^{\otimes j-1}} \otimes (\epsilon_{A^{\otimes i-1}} \otimes \phi(r) \otimes \epsilon_{A^{\otimes n-i}}) \otimes \eta_{A^{\otimes n-j}}$. Now if we take two matrices of this form that compose it is easy to see that F_A preserves their composition. \square

Secondly,

Lemma. *If R is a rig then an algebra over the PROP $\text{Mat}(R)$ is a bicommutative bialgebra with a map of rigs $R \rightarrow \text{End}(A)$.*

Proof. Suppose that $F : \text{Mat}(R) \rightarrow \mathcal{C}$ is a strict monoidal functor. We set $A = F(1)$, $\mu = F((11))$, and $\eta = F(0 : 0 \rightarrow 1)$. This makes (A, μ, η) into a commutative monoid:

Associativity follows from

$$(1 \ 1) \begin{pmatrix} 1 & 1 & 0 \\ 0 & 0 & 1 \end{pmatrix} = (1 \ 1 \ 1) = (1 \ 1) \begin{pmatrix} 0 & 1 & 1 \\ 1 & 0 & 0 \end{pmatrix},$$

commutativity from

$$(1 \ 1) \begin{pmatrix} 0 & 1 \\ 1 & 0 \end{pmatrix} = (1 \ 1),$$

and the unit axiom from

$$(1 \ 1) \begin{pmatrix} 0 \\ 1 \end{pmatrix} = (1) = (1 \ 1) \begin{pmatrix} 1 \\ 0 \end{pmatrix}.$$

The coalgebra structure is given by comultiplication $\Delta = F((11)^t)$ and $\epsilon = (0 : 1 \rightarrow 0)$ that this does define a coalgebra follows from what has gone before and the fact that the transpose map from $\text{Mat}(R)$ to itself is a contravariant functor that is self-inverse.

Next we must check that the algebra and coalgebra structures are compatible. It suffices to check that Δ is an algebra map and the counit respects the algebra structure. The first follows from the matrix equations

$$\begin{pmatrix} 1 \\ 1 \end{pmatrix} (1 \ 1) = \begin{pmatrix} 1 & 1 \\ 1 & 1 \end{pmatrix} = \begin{pmatrix} 1 & 0 & 1 & 0 \\ 0 & 1 & 0 & 1 \end{pmatrix} \begin{pmatrix} 1 & 0 \\ 1 & 0 \\ 0 & 1 \\ 0 & 1 \end{pmatrix},$$

and the latter from the fact that 0 is a terminal object in $\text{Mat}(R)$. \square

Finally we need to show

Lemma. *If R is a rig then a morphism between algebras F and G over $\text{Mat}(R)$ in \mathcal{C} is just a bialgebra map between $F(1)$ and $G(1)$ that commutes with the maps $R \rightarrow \text{End}(F(1))$ and $R \rightarrow \text{End}(G(1))$.*

Proof. Let's write A for $F(1)$ and B for $G(1)$

A natural transformation from F to G is a map $\theta : A \rightarrow B$ in \mathcal{C} such that if X is an $(m \times n)$ matrix with coefficients in R then $\theta^{\otimes n} F(X) = G(X) \theta^{\otimes m}$.

This condition for (1×1) -matrices says precisely that θ commutes with the maps $R \rightarrow \text{End}(A)$ and $R \rightarrow \text{End}(B)$. Then the condition for the matrix (11) implies that θ is an algebra map, and for (11)^t that it is a coalgebra map.

It now remains to show that a bialgebra map θ of the given form defines a natural transformation. Because F and G are enriched functors and θ^n is a bialgebra map $A^{\otimes n} \rightarrow B^{\otimes n}$ for every n it suffices to check the equation $\theta^{\otimes n} F(X) = G(X) \theta^{\otimes m}$ for $(m \times n)$ matrices with only one non-zero entry. This is a straightforward check. \square