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[6400]

#### ¶ 4. Proof of T 6017.

In this and the next paragraph the properties of the partial ordering  $P$  of D 6303, considered now as a notion of forcing, will be investigated. The convention on variables of ¶ 3 is maintained.

D 6400 Let  $\mathbb{B}$  be the algebra over  $P$ .

D 6401 Let  $\dot{x}$  be the element of  $V^{\mathbb{B}}$  defined by

$$\text{dom}(\dot{x}) = \{n \mid n < \omega\}$$

$$\dot{x}(n) = \sum^{\mathbb{B}} \{O_{(s, S)}^{\mathbb{B}} \mid \text{nes} \wedge (s, S) \in |P|\}.$$

As before, for  $(s, S) \in |P|$ ,  $\dot{o}_i \in \mathbb{L}^{\mathbb{B}}$ ,

D 6402  $(s, S) \Vdash \dot{o}_i \longleftrightarrow O_{(s, S)}^{\mathbb{B}} \leq [\dot{o}_i]^{\mathbb{B}}$ .

The next group of theorems, till T 6413, establish elementary properties of  $\mathbb{B}$  and  $\dot{x}$ .

T 6403 ZF  $\vdash [\dot{x} \leq \omega]^{\mathbb{B}} = 1$ .

Proof For  $i \in \text{dom}(\dot{x})$ ,  $[\dot{i} \in \omega]^{\mathbb{B}} = 1$ , so

$$[\dot{i} \in \dot{x}]^{\mathbb{B}} = \sum_{j \in \text{dom}(\dot{x})}^{\mathbb{B}} [\dot{i} = \dot{j}]^{\mathbb{B}} \cdot \dot{x}(\dot{j})$$

$$= \sum_{j \in \text{dom}(\dot{x})}^{\mathbb{B}} [\dot{j} \in \omega]^{\mathbb{B}} \cdot [\dot{i} = \dot{j}]^{\mathbb{B}} \cdot \dot{x}(\dot{j})$$

$$\leq \sum_{j \in \text{dom}(\dot{x})}^{\mathbb{B}} [\dot{j} \in \omega]^{\mathbb{B}} \cdot [\dot{i} = \dot{j}]^{\mathbb{B}} \cdot \dot{x}(\dot{j}) \text{ by identity axioms.}$$

$$\leq [\dot{i} \in \omega]^{\mathbb{B}};$$

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hence  $\{\forall i \in \omega \rightarrow \forall j \in \omega\}^B = \perp$  for all  $i, j$ ,

$$\therefore \{\forall x \in \omega\}^B = \perp. \quad \underline{\text{QED.}}$$

T6404 ZF  $\vdash \{\forall x \in \dot{x}\}^B = \dot{x}(x)$ .

Proof  $\{\forall x \in \dot{x}\}^B = \sum_{m < \omega} \{\forall n = m\}^B \cdot \dot{x}(m) = \dot{x}(x)$   
 $\therefore \{\forall n = m\}^B = \emptyset$  unless  $n = m$ .  
 $\underline{\text{QED.}}$

T6405 ZF  $\vdash \bigwedge \langle s, S \rangle / \forall n < \omega [$

$(n \in s \leftrightarrow \langle t, T \rangle \leq \langle s, S \rangle \vee \langle t', T' \rangle \leq \langle t, T \rangle \text{ net'})$

and  $(n \notin s \cup S \leftrightarrow \neg \langle t, T \rangle \leq \langle s, S \rangle \text{ (n } \in t))]$ .

Proof I)  $n \notin s$ , then  $V T \subseteq S \cap n \notin T$ ;

$\langle s, T \rangle \leq \langle s, S \rangle$ , and there is no

$\langle t', T' \rangle \leq \langle s, T \rangle$  s.t.  $n \in t'$ , as  $n \notin s \cup T$ .

If  $n \in s \cup S$ ,  $n \notin s$ , then set  $t = s \cup \{n\}$ ,

$T = \frac{S}{\{n\}}$ ; then  $\langle t, T \rangle \leq \langle s, S \rangle$ .  $\underline{\text{QED}}$

T6406 ZF  $\vdash \bigwedge \forall n < \omega / \forall \langle s, S \rangle :$

$$(i) \quad \dot{x}(x) = \bigcup \{O_{\langle s, S \rangle}^B \mid n \in s\}$$

$$(ii) \quad \langle s, S \rangle \Vdash n \in \dot{x} \iff n \in s$$

$$(iii) \quad \langle s, S \rangle \Vdash n \notin \dot{x} \iff n \notin s \cup S$$

$$(iv) \quad \langle s, S \rangle \Vdash \dot{s} \subseteq \dot{x} \subseteq s \cup S.$$

Proof. Immediate from the definition and TT 6404, 5.

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[6407]

Let  $M$  be a transitive model of ZF. I shall also consider these concepts contained in  $M$ . Write

D 6407  $\mathbb{P}^M = \langle |\mathbb{P}^M|, \leq^M \rangle$  for the set of  $M$   
satisfying  $\in M$  DD 6302, 3.

Thus  $|\mathbb{P}^M| = |\mathbb{P}| \cap M$ ;  $\leq^M$  is the  
restriction of  $\leq$  to  $|\mathbb{P}^M|$ , and it is true in  $M$  that  
 $\mathbb{P}^M$  is a partial ordering without minimal elements and  
the maximum element  $\langle 0, \omega \rangle$ .

D 6408 Write  $\mathbb{B}^M$  for the member of  $M$  which in  
 $M$  satisfies D 6400: that is, in  $M$   $\mathbb{B}^M$   
is the algebra over  $\mathbb{P}^M$ . Write  $M\mathbb{B}$  for the  
Booleanized universe constructed in  $M$  w.r.t. the  
algebra  $\mathbb{B}^M$ .

D 6409 Write  $\dot{x}^M$  for the element of  $M\mathbb{B}$  satisfying  
in  $M$  D 6401.

D 6410  $F \xrightarrow{M} x \iff$   $F$  is an  $M$ -generic filter  
on  $\mathbb{P}^M$ ,  $\tilde{F}$  the  $M$ -complete ultrafilter on  $\mathbb{B}^M$  it  
generates,  $\phi_{\tilde{F}} : M^{\mathbb{B}} \rightarrow M[F]$  the "collapsing"  
mapping of D 6207, and  $x = \phi_{\tilde{F}}(\dot{x}^M)$ .

D 6411  $x$  is  $\mathbb{P}$ -generic over  $M$   $\iff$   $\text{VF}(F \xrightarrow{M} x)$ .

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[6412]

T 6412 ZF  $\vdash$  Let  $F$  be an  $M$ -generic filter; and  $x \subseteq \omega$ .

The following are equivalent:

$$(a) F \xrightarrow{M} x$$

$$(b) x = \{n \mid (\exists^M \in x^M) \in F\}$$

$$(c) x = \bigcup \{s \mid \forall S \in M (S \subseteq \omega \wedge \langle s, S \rangle \in F)\}$$

$$(d) F = \{\langle s, S \rangle \in \mathbb{P}^M \mid s \subseteq x \subseteq s \cup S\}.$$

Proof: The equivalence of (a), (b) and (c) is immediate from the definitions. Note that  $s \subseteq x \subseteq s \cup S \rightarrow s \in x$ .

Suppose (c) (and therefore (a)) holds:

Then by (a),  $\langle s, S \rangle \in F \rightarrow s \subseteq x \subseteq s \cup S$ .  
(and T 6406(iv))

Conversely, if  $\langle s, S \rangle \in M$  and  $s \subseteq x \subseteq s \cup S$ ,

then  $\langle s, S \rangle$  is compatible with every element of  $F$ : for let  $\langle t, T \rangle \in F$ ; then  $t \subseteq x \subseteq t \cup T$ , so  $t \subseteq x$  and  $s \subseteq x$ , and therefore  $s \cup t = t \cup s$ ,  $t \setminus s \subseteq S$  and  $s \setminus t \subseteq T$ .

Then  $\langle s \cup t, S \cup T \rangle \leq^M \langle s, S \rangle$

and  $\langle s \cup t, S \cup T \rangle \leq^M \langle t, T \rangle$ .

Now the set  $\{\langle t, T \rangle \mid \langle t, T \rangle \leq^M \langle s, S \rangle \text{ or } \langle t, T \rangle$

is incompatible with  $\langle s, S \rangle (\in M)\}$

is in  $M$ , and is dense and  $\leq$ -closed.

$\therefore \exists \langle t, T \rangle \in F$  s.t.  $\langle t, T \rangle \leq \langle s, S \rangle$  (as it

has been already shown that no  $\langle t, T \rangle \in F$  is incompatible with  $\langle s, S \rangle$ ); hence  $\langle s, S \rangle \in F$  by D 6217 (ii); so (d) holds.

If (d) holds, let  $F \xrightarrow{M} y$ ; then for every  $t \in y$ ,  $t \subseteq x$ ; so  $y = x$ . QED.

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[6413]

T 6412 (c) and (d) together show that

T 6413 If  $F \xrightarrow{M} x$  then  $M[F] = M[x]$ .

T 6017 is now restated:

T 6017  $\text{ZF} \vdash$  Let  $M$  be a transitive model of  $\text{ZF} + \text{DC}$ ,  
 $x$   $\mathbb{P}$ -generic over  $M$  and  $y \subseteq x$ . Then  $y$   
is also  $\mathbb{P}$ -generic over  $M$ .

REMARKS (1) It is not asserted that for a general  
such  $M$  any such  $x$  exist.  
(2) It is not assumed that  $y \in M[x]$ .

The proof is in three steps: I shall give the  
first two as theorems in  $\text{ZF} + \text{DC}$  about  $\mathbb{P}$ . They will  
then hold in  $M$  about  $\mathbb{P}_M^M$ .

D 6414  $D =_{\text{df}} \{\Delta \subseteq |\mathbb{P}| \mid \Delta \text{ is dense and } \leq -\text{closed}\}$ .

D 6415 Let  $\Delta \in D$ ,  $\langle s, S \rangle \in |\mathbb{P}|$ .

$\langle s, S \rangle$  captures  $\Delta \longleftrightarrow \forall T \subseteq S \forall t \in T$  such  
that  $\langle s \cup t, \frac{S}{t} \rangle \in \Delta$ .

T 6416  $\text{ZF} + \text{DC} \vdash \forall \Delta \in D \forall \langle s, S \rangle \forall S' \subseteq S \langle s, S' \rangle \text{ captures } \Delta$ .

Proof Let  $\Delta$  be given.

(1) Define  $P_\Delta = \{X \subseteq \omega \mid \forall s \in X \langle s, \frac{X}{s} \rangle \in \Delta\}$ .

Then  $P_\Delta$  is CR+.

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For by T 6316, either  $P_\Delta$  is CR or for some  $\langle s, S \rangle$ ,  $P$  is CSF on  $2^{\langle s, S \rangle}$ . Let  $\langle s, S \rangle \in |P|$ . As  $\Delta$  is dense, there is an  $\langle s', S' \rangle \preceq \langle s, S \rangle$  such that  $\langle s', S' \rangle \in \Delta$ . But then  $2_{\omega}^{\langle s', S' \rangle} \subseteq P_\Delta$ , for if  $X \in 2_{\omega}^{\langle s', S' \rangle}$ , then  $s' \in X$  and  $\langle s', \frac{X}{s'} \rangle \preceq \langle s', S' \rangle$  and so  $\langle s', \frac{X}{s'} \rangle \in \Delta$  as  $\Delta$  is  $\preceq$ -closed; and  $\therefore X \in P_\Delta$ . It follows that  $P$  is CSF on no  $2^{\langle s, S \rangle}$ , and is therefore CR. So given  $\langle s, S \rangle \in |P|$  there is an  $S' \subseteq S$  such that either  $2_{\omega}^{\langle s, S' \rangle} \subseteq P_\Delta$  or  $2_{\omega}^{\langle s, S' \rangle} \cap P_\Delta = \emptyset$ ; but by the argument above, the second alternative is impossible, as there is an  $\langle s'', S'' \rangle \preceq \langle s, S' \rangle$ , with  $\langle s'', S'' \rangle \in \Delta$ ; and then  $s'' \cup S'' \in P_\Delta$ .

(1) is proved. Let now  $\langle s, S \rangle$  be given.

By (1) there is an  $S_0 \subseteq S$  such that  $2_{\omega}^{\langle s, S_0 \rangle} \subseteq P_\Delta$ . Let  $\langle s_n | n < \omega \rangle$  be the wonderful enumeration (of D 6313). Define a sequence of sets  $\langle S_i | i < \omega \rangle$  as follows:

$S_0$  is chosen s.t.  $s \in S_0$  and  $2_{\omega}^{\langle s, S_0 \rangle} \subseteq P_\Delta$ .  
at the  $n^{th}$  stage,

if  $s_n \notin S_n$ , set  $S_{n+1} = S_n$ ;

if  $s_n \subseteq S_n$  and there is an  $X \subseteq \frac{S_n}{s_n}$  such that  $\langle s \cup s_n, X \rangle \in \Delta$ , then pick such an  $X$  and set  $S_{n+1} = \frac{s \cup s_n}{S_n} \cup X$ ;

if there is no such  $X$ , set  $S_{n+1} = S_n$ .

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Then  $S_0 \supseteq S_1 \supseteq S_2 \supseteq \dots$ , and for all  $n$ ,  
 $s <_o S_n$ .

Let  $S' = \bigcap_{n<\omega} S_n$ . Then  $s <_o S'$

Exactly as in the proof of T 6314 it is seen  
that  $S'$  is infinite.

I assert that  $S'$  has the required properties. For  
let  $T \subseteq S'$ : then  $s <_o T$  and  $s \cup T \in 2_\infty^{(S, S_0)}$   
and so  $s \cup T \in P_\Delta$ . There is therefore a  $t' \in s \cup T$   
such that  $\langle t', \frac{T}{t'} \rangle \in \Delta$ . Now set  $t'' = s \cup t'$ :  
(note that either  $s \in t'$  or  $t' \in s$ ) then  $s \in t''$ ,  
and  $\langle t'', \frac{T}{t''} \rangle \preccurlyeq \langle t', \frac{T}{t'} \rangle \in \Delta$ , and so  
 $\langle t'', \frac{T}{t''} \rangle \in \Delta$ , as  $\Delta$  is  $\leq$ -closed. Write

$t'' = s \cup t$  where  $s <_o t$ . Then  $t = s_m$  some  $m$ ,  
and so  $\langle s \cup t, \frac{T}{s \cup t} \rangle \in \Delta$ ,  $S_{m+1}$  was of the  
form  $\overline{S_m} \cup X$ , where  $\langle s \cup t, X \rangle \in \Delta$ ; but then  
as  $\frac{S'}{s \cup t} \subseteq X$ ,  $\langle s \cup t, \frac{S'}{s \cup t} \rangle \in \Delta$ , as required.

( $\frac{S'}{s \cup t} = \frac{S'}{t} \rightsquigarrow s <_o S'$ , so

$$\langle s \cup t, \frac{S'}{s \cup t} \rangle = \langle s \cup t, \frac{S'}{t} \rangle .)$$

As in T 6314 it is seen that DC is  
enough for the whole proof.

Q.E.D.

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Let  $A = 2_\infty^\omega$ .

[6417]

T 6417  $ZF + DC \vdash \llbracket \forall \Delta \in \check{D} \forall S \in \check{A} (\dot{x} \subseteq S \text{ and } \dot{x} \in \dot{V}, \langle \dot{0}, S \rangle \text{ captures } \Delta) \rrbracket^B = 1$ .

T 6417 may be rephrased as

T 6418  $ZF + DC \vdash \llbracket \forall \Delta \in \check{D} \forall S \in \check{A} (\dot{x} \subseteq S \wedge \forall T \in \check{A} (T \subseteq S \rightarrow \exists t \in T \langle \langle t, \frac{S}{t} \rangle \in \Delta \rangle)) \rrbracket^B = 1$ ,

and has the corollary

T 6419  $ZF \vdash$  Let  $M$  be a transitive model of  $ZF + DC$ , and  $x$   $\mathbb{P}$ -generic over  $M$ . Then for any  $\Delta \in D^M$ , there is an  $S \in M$  such that  $x \subseteq S$  and in  $M$ ,  
 $\langle \dot{0}, S \rangle$  captures  $\Delta$ .

(Here  $D^M$  is in  $M$  and is the set that in  $M$  satisfies D 6414).

Derivation of T 6419 from T 6417

Let  $F \xrightarrow{M} x$ . T 6417 holds in  $M$  (replacing  $\dot{x}$  by  $\dot{x}^M$ ) and therefore as " $\phi_{\tilde{F}}(\dot{V}) = M$ " and  $\phi_{\tilde{F}}(\dot{x}^M) = x$ , the following is true in  $M[x]$ :

for any  $\Delta \in D^M$  there is an  $S \in M$  such that  $x \subseteq S$  and in  $M$ ,  $\langle \dot{0}, S \rangle$  captures  $\Delta$ .

As  $x \subseteq S \leftrightarrow x \subseteq S$  in  $M[x]$ , and the other quantifiers are relativised to  $M$ , the conclusion of T 6419 follows.

QED.

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[6420]

Proof of T 6417.

It suffices to show that

6420 if  $\Delta \in D$  and  $\langle s, S \rangle \in |P|$  then there is an  $\langle s', S' \rangle \preccurlyeq \langle s, S \rangle$  and an  $X \subseteq \omega$  such that

$$\langle s', S' \rangle \Vdash \dot{X} \ni \dot{x} \wedge \dot{\Delta} \subseteq \dot{X} (\dot{T} \in \dot{\Delta} \rightarrow \dot{V} \in \dot{T} (\dot{\langle t, \dot{X} \rangle} \in \dot{\Delta})),$$

as for such an  $\langle s', S' \rangle$ ,

$$\langle s', S' \rangle \Vdash \forall X \in \dot{\Delta} (\dot{x} \in \dot{X} \text{ and } \dot{V} \text{ captures } \dot{\Delta})$$

and so  $\{ \langle s, S \rangle \mid \langle s, S \rangle \Vdash \forall X \in \dot{\Delta} (\dot{x} \in \dot{X} \text{ and } \dot{V} \text{ captures } \dot{\Delta}) \}$

is dense in  $|P|$ , which shows that for each  $\Delta$ ,

that sentence has Boolean value  $\mathbb{I}$ .

Let then  $\Delta \in D$ ,  $\langle s, S \rangle \in |P|$ .

Enumerate the finite subsets of  $s$  as  $t_0, \dots, t_k$ .

Define a sequence  $S_0 \supseteq \dots \supseteq S_k$  as follows:

pick  $S_0 \subseteq S$  such that  $\langle t_0, S_0 \rangle$  captures  $\Delta$ :  
such an  $S_0$  exists by T 6416.

Pick  $S_1 \subseteq S_0$  such that  $\langle t_1, S_1 \rangle$  captures  $\Delta$ ;

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Pick  $S_k \subseteq S_{k-1}$  such that  $\langle t_k, S_k \rangle$  captures  $\Delta$ .

Remark that if  $\langle t, T \rangle$  captures  $\Delta$  and  $T' \subseteq T$ , then  $\langle t, T' \rangle$  captures  $\Delta$ . As  $S_k \subseteq S_i$  each  $i = 0, \dots, k$ ,

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[6421]

$\langle t_i, S_k \rangle$  captures  $\Delta$  for  $i = 0, 1, \dots, k$ .

I assert that  $\langle 0, s \cup S_k \rangle$  captures  $\Delta$ . For let  $T \subseteq s \cup S_k$ , and set  $t = T \cap s$ . Then  $t = t_i$  some  $i \leq k$ . Then  $T \in 2_\infty^{\langle t_i, S_k \rangle}$ , and as  $\langle t_i, S_k \rangle$  captures  $\Delta$ , there is a  $t' \in \frac{T}{t_i}$  such that  $\langle t_i \cup t', \frac{S_k}{t'} \rangle \in \Delta$ . But  $t_i \cup t' \in T$  and  $\frac{S_k}{t'} = \frac{S_k}{t_i \cup t'}$  ; so  $\forall t'' \in T \quad \langle t'', \frac{S_k}{t''} \rangle \in \Delta$ , as required.

Set  $X = s \cup S_k$ , and  $\langle s', S' \rangle = \langle s, S_k \rangle$ .

Then  $\langle s', S' \rangle \leq \langle s, S \rangle$

and  $\langle s', S' \rangle \Vdash \dot{x} \subseteq \dot{X}$  by T6406(iv).

As  $\langle 0, X \rangle$  captures  $\Delta$ ,

$\langle s', S' \rangle \Vdash \text{in } \dot{V}, \langle \dot{0}, \dot{X} \rangle \text{ captures } \dot{\Delta}$ .

QED

In the final step of the proof of T6017 I shall use a standard result which I give as T6423.

D6421 (ZF) A partial ordering  $\langle Q, R \rangle$  is a tree iff  
 (i)  $Q$  is countable and  $\forall q \in Q (\neg \exists q' \in Q (qRq'))$ .  
 (ii)  $\forall q \in Q \quad \{q' \mid qRq'\}$  is finite and linearly ordered by  $R$ .

REMARK None of the conventions about partial orderings assumed in

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[6422]

$\P \P 1$  and 2 are intended to apply in D 6421.

D 6422 (ZF) A tree  $\langle Q, R \rangle$  is well founded iff  
 there is no function  $g: \omega \rightarrow Q$  such that  
 $\forall i < \omega \ g(i+1) R g(i)$ ; that is, no  
 infinite  $R$ -descending paths.

T 6423 ZF  $\vdash$  A tree  $\langle Q, R \rangle$  is well-founded iff  
 there is a function  $f: Q \rightarrow \text{On}$  such  
 that  $\forall q, q' \in Q \ q R q' \rightarrow f(q) < f(q')$ .

Proof If there is an  $f: Q \rightarrow \text{On}$  such that  
 $\forall q, q' \in Q \ q R q' \rightarrow f(q) < f(q')$ , then  
 $\langle Q, R \rangle$  is well founded; for were  $g: \omega \rightarrow Q$   
 to be such that  $\forall i < \omega \ g(i+1) R g(i)$ , then  
 $\forall i < \omega \ f(g(i+1)) < f(g(i))$ , contradicting the  
 well-ordering of the ordinals.

and  $Q$  not empty. , suppose  $\langle Q, R \rangle$  is well-founded,

Let  $\Xi$  be the set of all functions  $\xi$  such

that

$$(a) \ \text{dom}(\xi) \subseteq Q$$

$$(b) \ \text{range}(\xi) \subseteq \text{On}.$$

$$(c) \ q \in \text{dom}(\xi) \wedge q' R q \rightarrow q' \in \text{dom}(\xi).$$

$$(d) \ \text{if } \neg \forall q' q' R q \text{ and } q \in \text{dom}(\xi),$$

$$\text{then } \xi(q) = 0.$$

$$(e) \ \text{if } \forall q' q' R q \text{ and } q \in \text{dom}(\xi),$$

$$\text{then } \xi(q) = \sup \{ \xi(q') + 1 \mid q' R q \}.$$

Then

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$$(1) \forall q \in Q \rightarrow \forall q' \in Q \quad q' R q.$$

Suppose not. Let  $\psi : \omega \longleftrightarrow Q$  be a well ordering of  $Q$ , which (such exist as  $Q$  is countable).

Define  $g(0) =$  the first element of  $Q$  in the well ordering  $\psi$ .

("the  $\psi$ -first").

$g(i+1) =$  the  $\psi$ -first  $q \in Q$  such that  
 $q R g(i)$ .

Then  $\forall i < \omega \quad g(i+1) R g(i)$ , contradicting the well foundedness of  $Q$ .

(2)  $\Xi$  is not empty.

For let  $q \in Q$  s.t.  $\neg \forall q' \in Q \quad q' R q$ .

Define  $\xi$  by  $\xi(q) = 0$ ,

$\xi$  undefined otherwise.

Then  $\xi \in \Xi$ .

(3) if  $\xi, \xi' \in \Xi$  and  $q \in \text{dom}(\xi) \cap \text{dom}(\xi')$

then  $\xi(q) = \xi'(q)$ :

for if  $\neg \forall q' \quad q' R q$ , then  $\xi(q) = \xi'(q) = 0$ ;  
 so if  $\xi(q) \neq \xi'(q)$ ,

$\forall q' (q' R q \text{ and } \xi(q') \neq \xi'(q'))$ ,  
 by (e), (c).

But then pick the  $\psi$ -first such  $q'$ : there is

a  $q'' R q' \quad \xi(q'') \neq \xi'(q'') \dots$ .

clearly there is then a map  $g : \omega \rightarrow Q$  s.t.

$\forall i \quad g(i+1) R g(i)$  - again contradicting wellfoundedness.

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Define  $f = \bigcup \{\xi \mid \xi \in \Xi\}$ .

Then by (3),  $f$  is a function,

$\text{dom}(f) \subseteq Q$ ,  $\text{range}(f) \subseteq \alpha$ .

$f(q)$  is defined iff  $\forall \xi \in \Xi$  s.t.

$\xi(q)$  is defined, and then  $f(q) = \xi(q)$ .

(4)  $f$  is everywhere defined.

For if  $q \in Q$  and  $\{q' \mid q' R q\} \subseteq \text{dom}(f)$ ,

then define  $\xi(q') = f(q')$   $q' R q$ ,

$$\xi(q) = \sup \{f(q') + 1 \mid q' R q\}.$$

Then  $\xi \in \Xi$ , and so  $q \in \text{dom}(f)$ .

So if  $\text{dom}(f) \neq Q$ ,

define  $g(0) = \text{the } \leftarrow\text{-first } q \text{ not in } \text{dom}(f)$ ;

then  $\forall q' q' R q$  and  $q' \notin \text{dom}(f)$ ;

let  $g(1) = \text{the } \leftarrow\text{-first such } q'$ .

$g(i+1) = \text{the } \leftarrow\text{-first } q \notin \text{dom}(f)$

s.t.  $q R g(i)$ ;

then  $\forall i g(i+1) R g(i)$  #.

So  $f: Q \rightarrow \alpha$  and  $q R q' \rightarrow f(q) < f(q')$ . QED

I shall now prove T 6017.

Let  $M$  be a transitive model of  $ZF + DC$ ,  $x$   $\mathbb{P}$ -generic over  $M$ , and  $y \subseteq x$ .

Define  $F_y = \{(s, s) \in M \mid s \in y \in \text{su } S\}$ .

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I assert that  $F_y$  is an  $M$ -generic filter on  $\mathbb{P}^M$ .

(1) Let  $\langle s, S \rangle \in F_y$  and  $\langle s', S' \rangle \in F_y$ .

Then  $s \leq y \leq s \cup S$ ,  $s' \leq y \leq s' \cup S'$ ,

and  $s \in_{\text{in } M} y$ ,  $s' \in_{\text{in } M} y$  so that

$s \in s'$  or  $s' \in s$ .

Suppose the first, without loss of generality. Then

$$s' \leq y \leq s' \cup \frac{S}{s'},$$

$$\Rightarrow y \leq s' \cup (S' \cap \frac{S}{s'}). \text{ Set } T = S' \cap \frac{S}{s'}.$$

Then  $\langle s', T \rangle \leq \langle s', S' \rangle$ ,  $\langle s', T \rangle \leq \langle s', \frac{S}{s'} \rangle \leq \langle s, S \rangle$ ,

and  $s' \leq y \leq s' \cup T$ , so  $\langle s', T \rangle \in F_y$ .

(2) Let  $\langle s, S \rangle \in F_y$ ,  $\langle s, S \rangle \leq \langle s', S' \rangle$ .

Then  $s' \leq s \leq y \leq s \cup S \leq s' \cup S'$ , so  $\langle s, S \rangle \in F_y$ .

(3) Let  $\Delta \in D^M$ : that is  $\Delta$  is in  $M$  a dense and  $\leq$ -closed subset of  $|\mathbb{P}^M|$ . By T6419 there is an  $S^\Delta \supseteq x$  such that  $S \in M$  and in  $M$ ,

$S^\Delta$  captures  $\Delta$ .

Let  $Q^\Delta = \{t \leq S^\Delta \mid \langle t, \frac{S^\Delta}{t} \rangle \notin \Delta\}$ .

$Q^\Delta \in M$ .

If  $Q^\Delta$  is empty, then  $\langle 0, S^\Delta \rangle \in \Delta$ ; but  $0 \leq y \leq x \leq S^\Delta = 0 \cup S^\Delta$ , so then  $\langle 0, S^\Delta \rangle \in F_y$ , as required.

Suppose that  $Q^\Delta$  is not empty. Define a relation

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$R^\Delta$  on  $Q^\Delta$  by

$t_1 R^\Delta t_2 \iff t_2 \in t_1 \text{ and } t_1 \neq t_2$ :

that is,  $t_1 R^\Delta t_2$  iff  $t_1$  is a proper end extension of  $t_2$ .  $R^\Delta \in M$ . Set  $T^\Delta = \langle Q^\Delta, R^\Delta \rangle$ .

T 6424 In  $M$ ,  $T^\Delta$  is a well-founded tree.

Proof. (A)  $R^\Delta$  is transitive, irreflexive, and for all  $s \in Q$ ,

$$s R^\Delta t_1 \& s R^\Delta t_2 \rightarrow t_1 \in s \text{ and } t_2 \in s$$

$$\rightarrow t_1 \in t_2 \text{ or } t_2 \in t_1$$

$$\rightarrow t_1 R^\Delta t_2, t_1 = t_2 \text{ or } t_2 R^\Delta t_1.$$

So as  $S$  has only finitely many initial segments,

$\{t \mid s R^\Delta t\}$  is finite, and linearly ordered by  $R^\Delta$ .

(B)  $Q^\Delta$  is countable, being a subset of the set of finite subsets of  $\omega$ .

(C) Suppose  $\langle t_i \mid i < \omega \rangle \in M$  is a sequence such that  $\forall i \exists t_{i+1} R^\Delta t_i$ : so that in particular  $t_i \in Q^\Delta$  each  $i$ , and  $i < j \rightarrow t_i \in t_j$ .

Set  $T = \bigcup_{i < \omega} t_i$ .

Then  $T \subseteq S$ , and as  $S$  captures  $\Delta$ , there is a  $t \in T$  such that  $\langle t, \frac{S}{t} \rangle \in \Delta$ . But then for some  $i$ ,  $t \in t_i$ ; and  $\langle t_i, \frac{S}{t_i} \rangle \leq \langle t, \frac{S}{t} \rangle$ ; as  $\Delta$  is  $\leq$ -closed,  $\langle t_i, \frac{S}{t_i} \rangle \in \Delta$ , and so  $t_i \notin Q^\Delta$ .  $\#\$ .

QED

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[6425]

Now  $M$  is a model of  $ZF$ , and therefore by T6423,  
there is an  $f: Q^\Delta \rightarrow \text{On}^M$ , ( $f \in M$ ) such that

$$\forall q, q' \in Q^\Delta \quad q R^\Delta q' \rightarrow f(q) < f(q').$$

Now  $M$  is transitive, and so the ordinals of  $M$  are an initial segment of the ordinals in  $V$ , so in  $V$ ,  $f$  is a function  $f: Q^\Delta \rightarrow \text{On}$ , and

$$\forall q, q' \in Q^\Delta \quad q R^\Delta q' \rightarrow f(q) < f(q').$$

Applying T6423 (in  $V$ ),

T6425  $T^\Delta$  is a well founded tree.

$y \subseteq x \subseteq S^\Delta$ ; so by T6425, there is an  $s \in y$  such that  $\langle s, \frac{S}{s} \rangle \in \Delta$  (otherwise  $\{s \mid s \in y\}$  would form an infinite  $R^\Delta$ -descending path in  $Q^\Delta$ , contradicting T 6425.) But then  $s \subseteq y \subseteq s \cup \frac{S}{s}$ , so  $\langle s, \frac{S}{s} \rangle \in F_y \cap \Delta$ .

I have shown therefore that

$$\forall \Delta \in D^M \quad F_y \cap \Delta \neq \emptyset.$$

That, and points (1), (2) on page 75 show that  $F_y$  is indeed an  $M$ -generic filter on  $P^M$ .

By T6412,  $F_y \xrightarrow{M} y$ , and so  $y$  is  $P$ -generic over  $M$ .

The proof of T6017 is complete.

(78)

[6426]

D 6426 Let  $ZF_n$  denote the conjunction of the first  $n$  axioms of  $ZF$  in some fixed recursive enumeration such that the axiom of extensionality is first, the axiom of foundations the second, and the axiom of infinity is third.

In defining  $\langle |P|, \leq \rangle$ ,  $B$ ,  $V_\alpha^B$ , the functions  $[x \in y]$ ,  $[x = y]$  (which are not sets), and  $D$ , the set of dense and  $\leq$ -closed subsets of  $|P|$ , certain axioms of  $ZF$  have been used to ensure that there are sets satisfying the various definitions: only finitely many of the axioms have been used, say a subset of the first  $m_0$ .

For  $M$  a transitive model of  $ZF_{m_0} + DC$ , " $x$  is  $P$ -generic over  $M$ " may be defined as before: that is, that there is an  $F \subseteq |P^M|$  such that for all  $\Delta \in D^M$ ,  $F \cap \Delta \neq \emptyset$ , and  $F \xrightarrow{M} x$ .

If  $n \geq m_0$  and  $M$  is a transitive model of  $ZF_n + DC$ , and  $x$  is  $P$ -generic over  $M$ , then  $M[x]$  will satisfy at least those axioms  $\Omega$  of  $ZF$  for which  $[\Omega]$  can be defined in  $M$  and proved to equal 11.

T 6427  $ZF \vdash$  There is an  $n \geq m_0$  such that if  $M$  is a transitive model of  $ZF_n + DC$ ,  $x$   $P$ -generic over  $M$  and  $y \in x$ , then  $y$  is also  $P$ -generic over  $M$ .

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Proof. By examining the proof of T6017, which used the cuts in  $M$  of only a finite number of axioms of  $ZF$ . That is, pick  $n$  so large that the definitions of the following statements are meaningful, and such that the statement are theorems of  $ZF_n + DC$ , when written with all variables bound and defined terms replaced by their defining clauses written out in full:

T6313 and its corollaries, TT6315, 16; T6321, T6412, T6415, T6418 of which T6419 will then be a corollary; T6423; and T6424, (with  $\Delta$  a bound variable).

QED.

D6428 Let  $(\underline{n})$  be the least such  $n$ .

T6427 will not be used in the proof of T6001, but in applying the notion of a  $P$ -generic real to establish Silver's theorem T6013.

I conclude the paragraph with three further observations about  $P$ : the first is due to Jensen, and will simplify some later proofs.

T6429 (Jensen)  $ZF + DC \vdash$  Let  $\dot{\alpha} \in L^B$ , and  $\langle s, S \rangle \in |P|$ . ( $B$  is the algebra over  $P$ ). Then there is an  $S' \subseteq S$  such that  
 $\langle s, S' \rangle \Vdash \dot{\alpha} \text{ or } \langle s, S' \rangle \Vdash \neg \dot{\alpha}$ .

Proof: (not Jensen's) The set  $\Delta =_{df} \{ \langle t, T \rangle \mid \langle t, T \rangle \Vdash \dot{\alpha} \text{ or } \langle t, T \rangle \Vdash \neg \dot{\alpha} \}$

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[6430]

is dense and  $\leq$ -closed, by the basic property of forcing F 1, page 21. By T 6416 there is an  $S'' \subseteq S$  such that  $\langle s, S'' \rangle$  captures  $\Delta$ . Consider the family

$$P = \{ X \in 2_{\omega}^{(s, S'')} \mid \forall t \text{ in } X \langle t, \frac{S''}{t} \rangle \Vdash \emptyset \}.$$

Let  $\langle t, T \rangle \preccurlyeq \langle s, S'' \rangle$ ; there is a  $t'$  in  $T$  such that  $\langle t \cup t', \frac{S''}{t'} \rangle \in \Delta$ , as  $\langle s, S'' \rangle$  captures  $\Delta$ .

Then  $2_{\omega}^{(t \cup t', \frac{S''}{t'})} \subseteq P$  or  $\subseteq S(\omega) \setminus P$ ; so  $P$  is not CSF on any  $\langle t, T \rangle$  with  $\langle t, T \rangle \preccurlyeq \langle s, S'' \rangle$ , and therefore  $P$  is CR on  $2^{(s, S'')}$ , by T 6316 applied to  $2^{(s, S'')}$  rather than  $2^{(0, \omega)}$ . So there is an  $S' \subseteq S''$  such that either  $2_{\omega}^{(s, S')} \subseteq P$  or  $2_{\omega}^{(s, S')} \cap P = \emptyset$ .

Suppose the first. Then  $\langle t, T \rangle \preccurlyeq \langle s, S' \rangle$  does

$$\langle t, T \rangle \Vdash \neg \emptyset,$$

$t \cup T \notin P$ , (as  $S' \subseteq S''$ ). So therefore  $\langle s, S' \rangle \Vdash \emptyset$ .

If the second holds,  $\langle s, S' \rangle \Vdash \neg \emptyset$ , by a similar argument.

Q.E.D.

T 6430 2F+ Let  $B$  be the algebra over  $P$ . Then  
 $B$  is homogeneous.

Proof: By T 6126, it is enough to show that for all  $p \in |P|$ ,  $B|_0 \cong B$ .

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[6431]

$\{O_q \mid q \leq p\}$  is dense in  $\mathbb{B} \setminus O_p$ , and so

$\mathbb{B} \setminus O_p \cong$  the algebraic  $\langle \{q \mid q \leq p\}, \leq \rangle$ :

but let  $p = \langle s, S \rangle$ , in the homeomorphism (D 6305)

$$h: 2^{\langle s, S \rangle} \xrightarrow{\sim} 2^\omega.$$

Then  $h$  is an isomorphism w.r.t.  $\leq$ ,

and so  $\mathbb{B} \setminus O_p \cong \mathbb{B}$ .

Q.E.D.

T6431.  $ZF \vdash$  Let  $M$  be a transitive model of  $ZF + DC$ .

Suppose there is a real  $P$ -generic over  $M$ . Then for all  $\langle s, S \rangle \in |P^M|$  there is an  $M$ -complete filter on  $P^M$  containing  $\langle s, S \rangle$ .

Proof: Let  $x$  be  $P$ -generic over  $M$ , and let  $F_x \rightarrow_M x$ , and  $\tilde{F}_x$  the  $M$ -complete ultrafilter on  $\mathbb{B}$  generated by  $F_x$ . Let  $\langle s, S \rangle \in |P^M|$ ; pick  $b \in F_x$ ,  $b \neq \mathbb{1}$ .

Set  $b' = O_{\langle s, S \rangle}$ . By T6430, there is an automorphism  $\phi \in \mathbb{B}$ , with  $\phi(b) = b'$ . Let  $\tilde{F} = \{\phi(c) \mid c \in \tilde{F}_x\}$ . I assert that  $\tilde{F}$  is an  $M$ -complete ultrafilter on  $\mathbb{B}$ :

(i) as  $\phi$  is an automorphism,  $\tilde{F}$  is an ultrafilter.

(ii) let  $\mathcal{X} \subseteq \tilde{F}$ ; then  $\{c \mid \phi(c) \in \mathcal{X}\} \subseteq \tilde{F}_x$  and is in  $M$  as  $\phi$  and  $\mathcal{X}$  are both in  $M$ , and so

$\text{Tr} \{\langle c \mid \phi(c) \in \mathcal{X} \rangle \in \tilde{F}_x\} \in \tilde{F}$ ; so  $\text{Tr} \mathcal{X} = \phi(\text{Tr} \{\langle c \mid \phi(c) \in \mathcal{X} \rangle\}) \in \tilde{F}$ .

Therefore  $\{\langle t, T \rangle \mid O_{\langle t, T \rangle} \in \tilde{F}\}$  is an  $M$ -generic filter on  $P^M$  containing  $\langle s, S \rangle$ . Q.E.D.