

MATHEMATICAL TRIPOS PART II (2005–06)

Coding and Cryptography - Example Sheet 4 of 4

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- 46) What is a linear feedback shift register? Show that, subject to a suitable non-degeneracy condition, any output stream x_0, x_1, x_2, \dots produced is purely periodic, i.e. there exists r such that $x_{n+r} = x_n$ for all $n \geq 0$.
- 47) A linear feedback shift register was used to generate the following stream

110001110001...

Recover the feedback polynomial by the Berlekamp-Massey method. [The LFSR has length 4 but you should work through the trials for length d for $1 \leq d \leq 4$.]

- 48) We work with streams of symbols in \mathbb{F}_2 . I have a secret sequence k_1, k_2, \dots and a message p_1, p_2, \dots, p_N . I transmit $p_1 + k_1, p_2 + k_2, \dots, p_N + k_N$ and then, by error, transmit $p_1 + k_2, p_2 + k_3, \dots, p_N + k_{N+1}$. Assuming that you know this and that my message makes sense, how would you go about finding my message? Can you now decipher other messages sent using the same part of my secret sequence?
- 49) One of the most confidential German codes (called FISH by the British) involved a complex mechanism which the British found could be simulated by two loops of paper tape of length 1501 and 1497. If $k_n = x_n + y_n$ where x_n is a stream of period 1501 and y_n is stream of period 1497, what is the longest possible period of k_n ? How many consecutive values of k_n do you need to specify the sequence completely?
- 50) A non-linear feedback register of length 4 has defining relation

$$x_{n+1} = x_n x_{n-1} + x_{n-3}.$$

Show that the state space contains 4 cycles of lengths 1, 2, 4 and 9.

- 51) The modulus $N = 713$ is used for the Rabin-Williams code. The ciphertext received is $c = 289$. Determine all possible plaintexts.
- 52) I announce that I shall be using the Rabin-Williams scheme with modulus N . My agent in X'Dofro sends me a message m (with $1 \leq m \leq N-1$) encoded in the requisite form. Unfortunately, my cat eats the piece of paper on which the prime factors of N are recorded so I am unable to decipher it. I therefore find a new pair of primes and announce that I shall be using the Rabin Williams scheme with modulus $N' > N$. My agent now recodes the message and sends it to me again.

The dreaded SNDO of X'Dofro intercept both code messages. Show that they can find m . Can they decipher any other messages sent to me using only one of the coding schemes?

- 53) (i) A user of RSA accidentally chooses a large prime for his modulus N . Explain why this system is not secure.
- (ii) A popular choice for the RSA encryption exponent is $e = 65537$. Using this exponent how many multiplications are required to encrypt a message?
- (iii) Why might it be a bad idea to use an RSA modulus $N = pq$ with $|p - q|$ small?

- 54) Alice and Bob are issued with RSA public keys (N, e_1) and (N, e_2) , and corresponding private keys (N, d_1) and (N, d_2) .
- The same message m is sent to both Alice and Bob. Assuming e_1 and e_2 are coprime, how can we recover m from the intercepted ciphertexts c_1 and c_2 ?
 - How can Alice read messages sent to Bob?
- 55) Alice and Bob carry out a Diffie-Hellman key exchange with $p = 29$ and $g = 2$. If the numbers exchanged by Alice and Bob are 9 and 12, what is their secret key?
- 56) Extend the Diffie-Hellman key exchange system to cover three participants in a way that is likely to be as secure as the two party scheme.
- Extend the system to n parties in such a way that they can compute their common secret key in at most $n^2 - n$ communications. [The numbers p and g of our original Diffie-Hellman system are known by everybody in advance.]
- 57) Give an example of a homomorphism attack on an RSA code. Show in reasonable detail that the el Gamal signature scheme (even without the use of a hash function) defeats it.
- 58) Suppose we drop the requirement $1 \leq r \leq p - 1$ from the el Gamal signature scheme. How might we then be able to forge new signatures from old? [Hint: Use the Chinese Remainder Theorem for the coprime moduli p and $p - 1$.]
- 59) Criticise the following authentication procedure. Alice chooses public key N for the Rabin-Williams code. To be sure we are in communication with Alice we send her a “random item” $r = m^2 \pmod{N}$. On receiving r , Alice proceeds to decode using her knowledge of the factorisation of N , and finds a square root m_1 or r . She returns m_1 to us and we check that $r = m_1^2 \pmod{N}$. [You should think about what happens when many mutually distrusting parties communicate with Alice in this way.]
- 60) Let K be the finite field with 2^d elements. We recall that K^* is a cyclic group, generated by α say. Let $T : K \rightarrow \mathbb{F}_2$ be any non-zero \mathbb{F}_2 -linear map.
- Show that the \mathbb{F}_2 -bilinear form $K \times K \rightarrow \mathbb{F}_2 ; (x, y) \mapsto T(xy)$ is non-degenerate (i.e. $T(xy) = 0$ for all $y \in K$ implies $x = 0$).
 - Show that the sequence $x_n = T(\alpha^n)$ is the output from a LFSR of length d .
 - The period of $(x_n)_{n \geq 0}$ is the least integer $r \geq 1$ such that $x_{n+r} = x_n$ for all sufficiently large n . Show that the sequence in (ii) has period $2^d - 1$.