

# HIGHER DIMENSIONAL CATEGORIES

- Categories to start higher categories
- Strict monoidal categories I
- $\Delta$  and simplicial sets
- Strict monoidal categories II
- Bicategories and the pentagon
- Higher dimensional category theory

# **Higher Dimensional Categories**

**Martin Hyland**

According to Eilenberg and Mac Lane, founders of Category Theory, categories were needed to explain functors and functors were needed to explain natural transformations. Thus from the beginning it was clear that the study of 1-dimensional categories required a 2-dimensional structure, a 2-category; and so on. In the simplest versions of the resulting higher dimensional categories, one takes composition to be strictly associative. I shall outline the theory and explain some results.

Unfortunately, for compelling applications the associativity of composition fails, and one needs a theory which treats ever more subtle notions of equivalence. One approach to this borrows ideas from homotopy theory. I shall try both to explain some concrete motivations and to convey the flavour of the subject as it appears to be developing.

CATEGORIES

TO

STRICT HIGHER

CATEGORIES

# CATEGORIES

A category  $\mathbb{C}$  has

objects  
points  
 $Q$ -cells

$A, B, C, \dots$

arrows  
maps morphisms  
 $1$ -cells

$A \xrightarrow{f} B$

identities

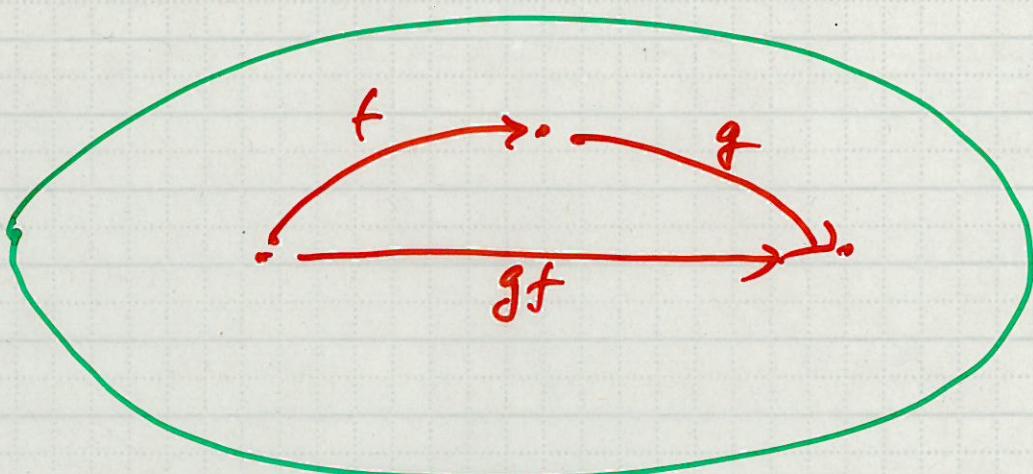
$A \xrightarrow{\gamma_A} A$

composition

$$\begin{array}{c} \underline{A \xrightarrow{f} B \quad B \xrightarrow{g} C} \\ A \xrightarrow{gf} C \end{array}$$

with evident axioms.

## PICTURE



Given  $C$ , we have  $C^{\text{op}}$  with  
same object but direction of arrows  
reversed.

$$(C^{\text{op}})^{\text{op}} = C$$

# USUAL EXAMPLES

	OBJECTS	ARROWS
•	Sets	Functions
•	Sets	Relations
•	Groups	Homomorphisms
•	Topological spaces	Continuous maps
•	Topological groups	Continuous homomorphisms
•	Functors ( $\mathcal{C}, \mathcal{D}$ )	Natural transformations

## SMALL CATEGORIES

A partially ordered set  $(P, \leq)$

gives a category with

objects : points of  $P$

maps : unique such  
 $p \rightarrow q$  when  $p \leq q$

A group  $G$  gives a category

with

objects : just one  $= *$  say

maps : elements of  $G$ ,

$g : * \rightarrow *$

(Composition = Multiplication).

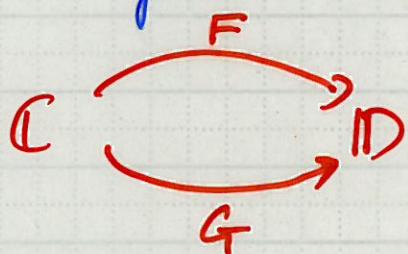
# FUNCTOR & NATURAL TRANSFORMATIONS

A structure preserving map

$$F: \mathcal{C} \rightarrow \mathcal{D}$$

between categories is a functor.

Given two functors



a natural transformation

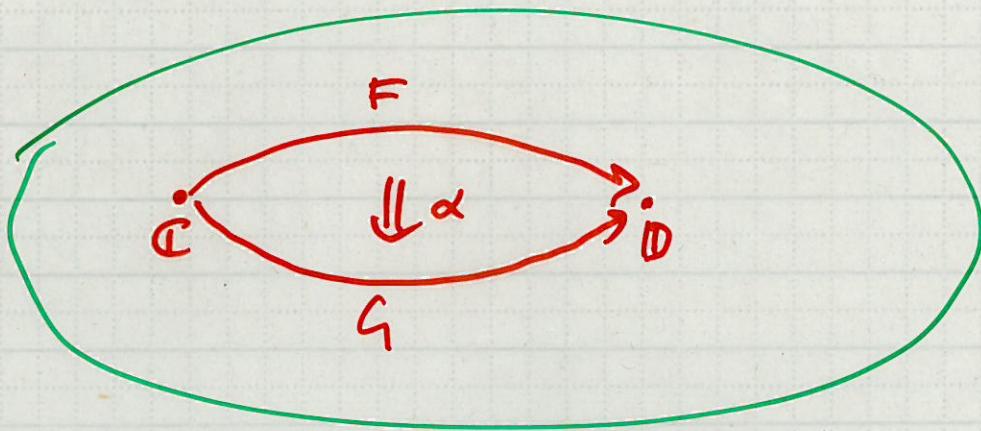
$$\alpha: F \rightarrow G$$

is given by  $\alpha_c : F(c) \rightarrow G(c)$  in  $\mathcal{D}$

with  $G(u) \cdot \alpha_c = \alpha_{D(u)} \cdot F(u)$  all  
 $u: C \rightarrow D$  in  $\mathcal{C}$

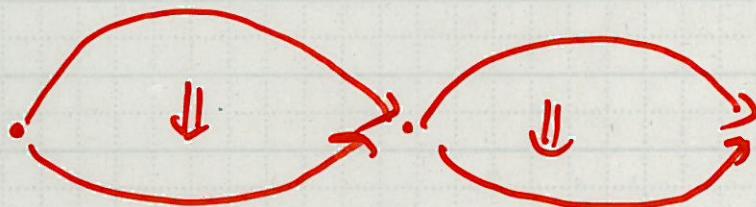
(forget direction issue!)

# PICTURES

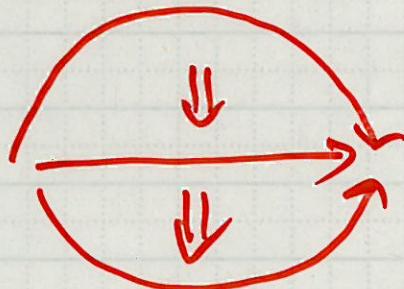


with two compositions

\*<sub>0</sub>



\*<sub>1</sub>



## 2 - CATEGORIES

A 2-category  $\mathcal{K}$  has

0-cells  $A, B, C, \dots$

1-cells  $A \xrightarrow{f} B, B \xrightarrow{g} C, \dots$

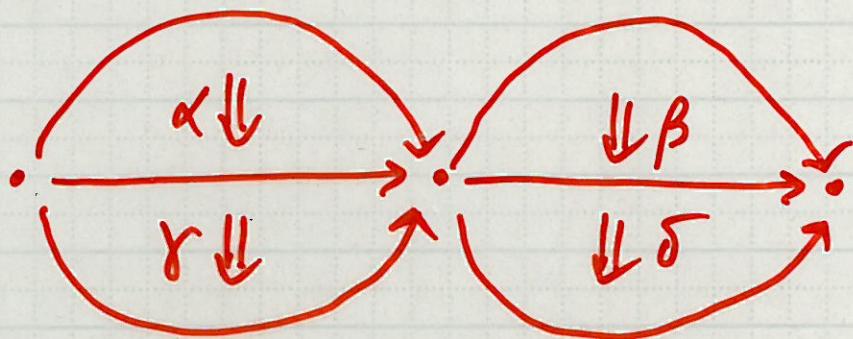
2-cells

The diagram shows two objects, A and B, connected by a horizontal arrow labeled f. A curved arrow labeled alpha points from A down to B, indicating a 2-cell between them.

with identities and compositions

( $\alpha$  two compositions of  
2-cells suffice)

## MIDDLE 4 INTERCHANGE

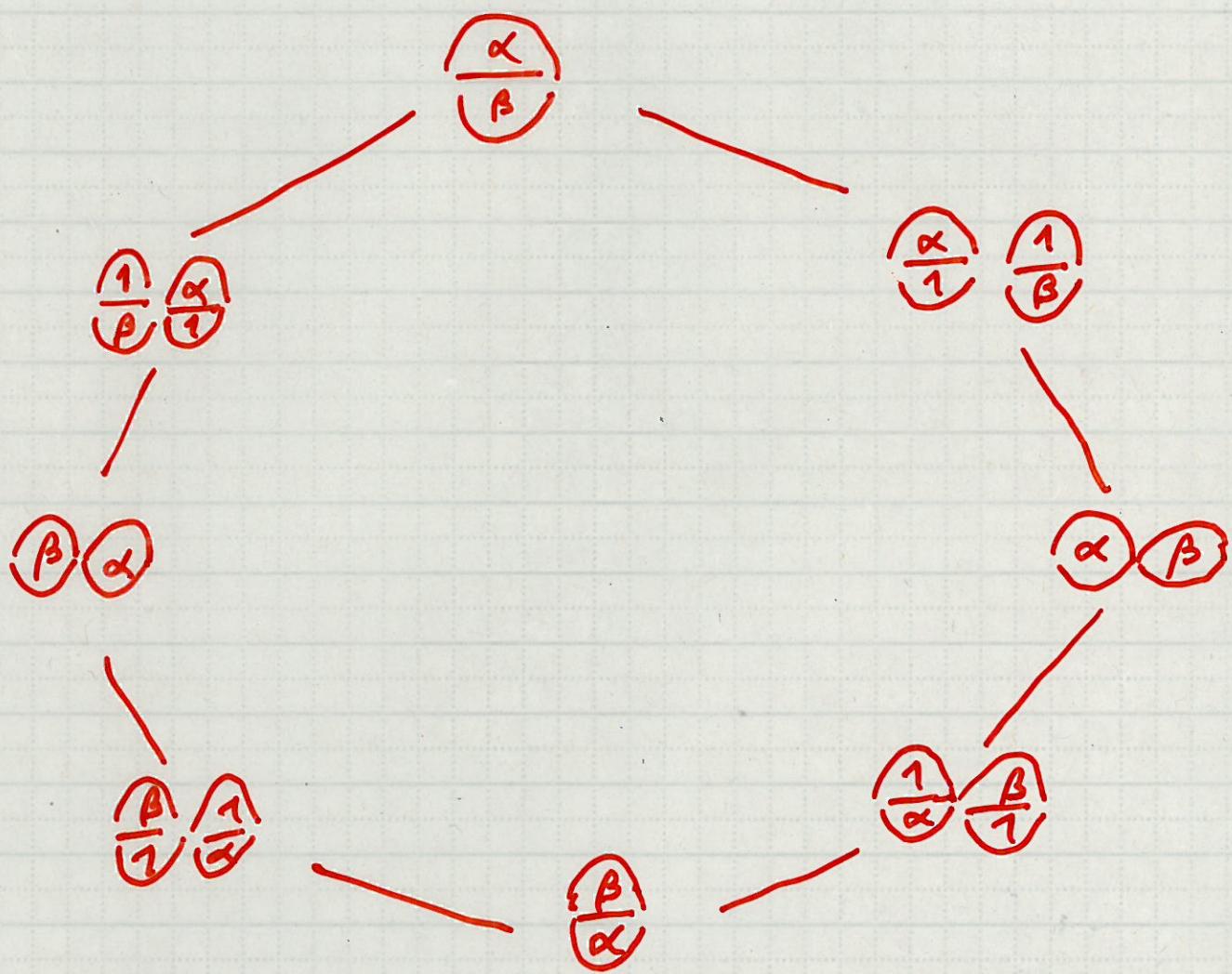


$$(\delta *_{\circ} \gamma) *_1 (\beta *_{\circ} \alpha) = (\delta *_1 \beta) *_{\circ} (\gamma *_1 \alpha)$$

This says the compositions  
commute

## ECKMANN - HILTON

Commuting compositions with 1 are equal,



and the composition is commutative.

## TWO CONSEQUENCES OF E-H ARGUMENT

Let  $A$  be a 0-cell (object)  
in a 2-category (or bicategory)  
 $\mathcal{K}$ . Then

$\text{End}(\mathbb{1}_A)$   
is a commutative monoid.

The second homotopy group  
 $\pi_2(X, x_0)$  of a pointed space is  
commutative (as are the higher  
homotopy groups).

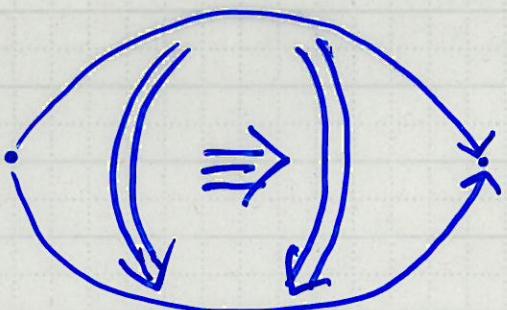
## STRICT HIGHER CATEGORIES

$$\dots \rightarrow C_3 \xrightarrow[s]{t} C_2 \xrightarrow[s]{t} C_1 \xrightarrow[s]{t} C_0$$

with

$$ss = st \quad (\text{2nd order } s)$$

$$ts = tt \quad (\text{2nd order } t)$$



globular  
sets

with identities and compositions  
at all levels.

STRICT  
MONOIDAL  
CATEGORIES

I

## (STRICT) MONOIDAL CATEGORIES

These are categories  $\mathbb{C}$  equipped with

$$I \in \mathbb{C} \quad \mathbb{C} \times \mathbb{C} \xrightarrow{\otimes} \mathbb{C}$$

such that

$$I \otimes A = A = A \otimes I$$

$$A \otimes (B \otimes C) = (A \otimes B) \otimes C$$

(These do not occur in the usual examples: here we have  $\cong$  for  $=$ .)

These are strict 2-categories with 1 object.

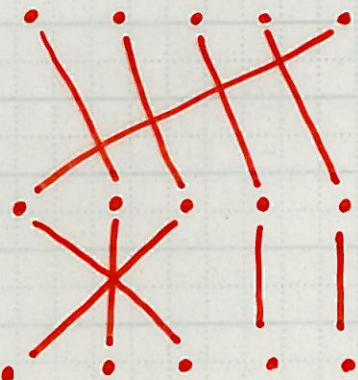
# SYMMETRIC GROUPS

objects : natural numbers  $n$

maps : bijections  $n \rightarrow n$

$\otimes$  : num (concatenation)

## PICTURE



$$"(13)(12345) = (12)(345)"$$

Generated as a monoidal category by



## BRAID GROUPS

objects

$n, \dots$

maps

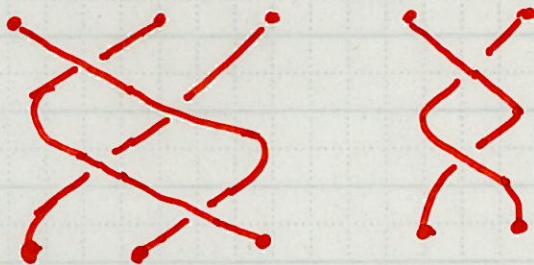
braids

$n \rightarrow n$

tensor

+

PICTURE



Generated by  $\times$  and its inverse  $\times$

as a monoidal category.

## THE BRAID RELATION



$$(S_{12} \otimes 1) (1 \otimes S_{23}) (S_{12} \otimes 1) = (1 \otimes S_{23}) (S_{12} \otimes 1) (1 \otimes S_{23})$$

$$R_{12} R_{23} R_{12} = R_{23} R_{12} R_{23}$$

cf Yang - Baxter equation

For the symmetric group add

$$S^2 = 1.$$

$\Delta$  AND  
SIMPLICIAL  
SETS

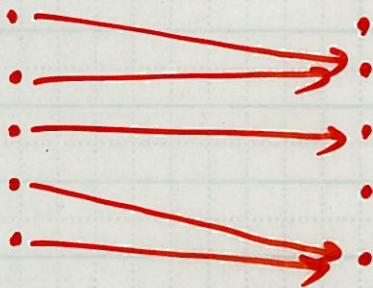
## FREE MONOID

$\Delta^+$  has

objects  $n = \Delta[n-1]$  (as an ordered set)

maps  $n \rightarrow m$  are order preserving maps

### PICTURE



Generated as a monoidal category by

$$\cdot \xrightarrow{\quad} \cdot \text{ and } \longrightarrow \cdot$$

## Axioms for $\Delta_+$

$$\begin{array}{c} \rightarrow \\ \vdots \\ \rightarrow \end{array} : \begin{array}{c} \rightarrow \\ \vdots \\ \rightarrow \end{array} = \begin{array}{c} \rightarrow \\ \vdots \\ \rightarrow \end{array}$$

$$\begin{array}{c} \rightarrow \\ \vdots \\ \rightarrow \end{array} : \begin{array}{c} \rightarrow \\ \vdots \\ \rightarrow \end{array} = \begin{array}{c} \rightarrow \\ \vdots \\ \rightarrow \end{array}$$

$$\begin{array}{c} \vdots \\ \rightarrow \\ \vdots \end{array} : \begin{array}{c} \rightarrow \\ \vdots \\ \rightarrow \end{array} = \begin{array}{c} \rightarrow \\ \vdots \\ \rightarrow \end{array}$$

The unit and associativity

laws for a monoid.

## SIMPLICIAL SETS

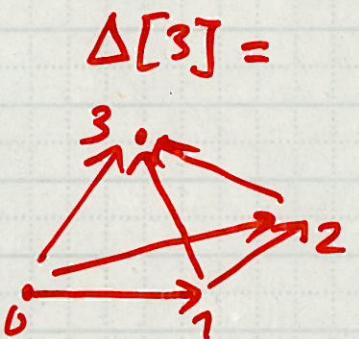
Let  $\Delta$  be as  $\Delta_+$  but omitting  
 $0 = \Delta[-1]$ .

So  $\Delta$  has objects

$$\Delta[0] = \bullet_0$$

$$\Delta[1] = \bullet_0 \xrightarrow{\quad} \bullet_1$$

$$\Delta[2] = \bullet_0 \xrightarrow{\quad} \bullet_1 \xrightarrow{\quad} \bullet_2$$



(i.e. formal simplices)

The category of simplicial sets is  
 the functor  $[\Delta^{\text{op}}, \text{Sets}]$  i.e presheaves  
 category on  $\Delta$

[The consequence of formally gluing  
 simplices together.]

$$X \in [\Delta^{\text{op}}, \text{Set}] : X_n = X(\Delta[n])$$

## THE NERVE OF A CATEGORY

Regard the ordered sets  $\Delta[n]$  as categories:  
order preserving maps are functors.

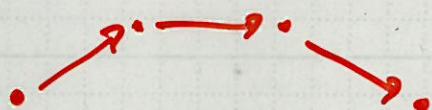
So if  $C$  is a category then

$\Delta(n) \longmapsto [\Delta[n], C]$  gives  
a simplicial set  $N(C)$ .

$N(C)_n = \text{composable } n\text{-paths}$   
 $\text{in } C$

E.g.

$N(C)_3$  is the set of diagrams



in  $C$

( $N(-)$  has a left adjoint, the fundamental category of a simplicial set.)

STRICT  
MONOIDAL  
CATEGORIES

II

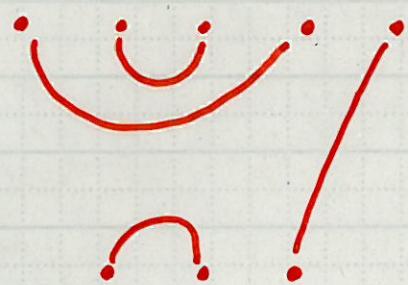
## TEMPERLEY - LIEB CATEGORY

The free monoidal category on a self dual object of dimension 1.

objects =  $n, \dots$

maps  $n \rightarrow m$  are chord diagrams

PICTURE



Generated as a monoidal category by



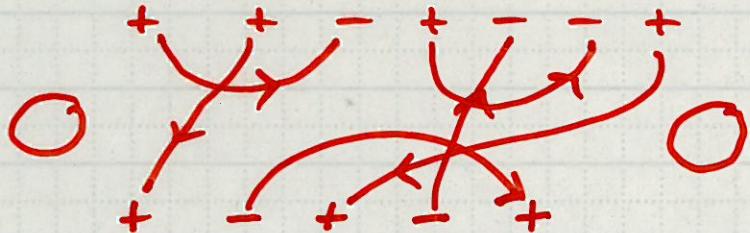
$$\text{Y} = \text{I} = \text{N} \quad \text{O} = \text{*}$$

## DUALS

Free category on object with involutive duals.

objects	sequences of signed points (closed zero dimensional manifolds)
maps	bijections respecting signs plus dummy circles (1-dimensional manifolds with appropriate $\partial$ )

## PICTURE



# TANGLES

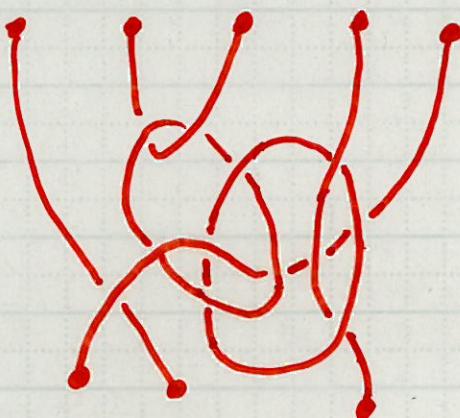
objects

$n, m, \dots$

maps

$n \rightarrow m$  tangles with  
 $n, m$  endpoints

PICTURE



Free braided monoidal on a self dual.

## APPLICATION

We have "structure preserving functors" to "naturally occurring" braided monoidal categories. (Determined by where the object  $\otimes$  goes.)

Eg to categories of representations of quantum groups, level  $k$  representations of loop groups.

A knot or link  $\longmapsto$  Map  $I$  to  $I$

and

$$\text{End}(I) = \mathbb{C}[q, q^{-1}]$$

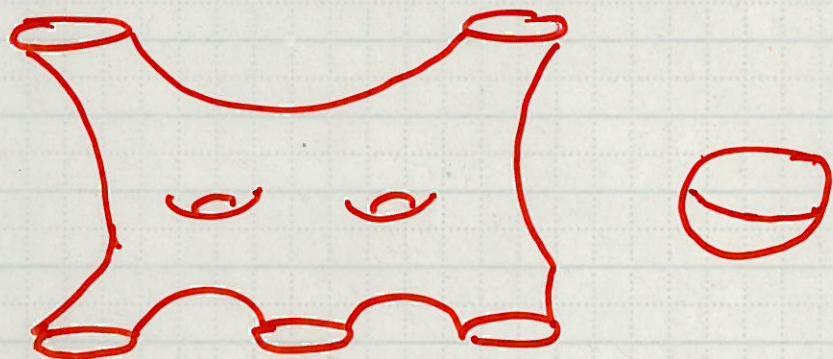
So knot invariants.

# NAIVE STRING THEORY

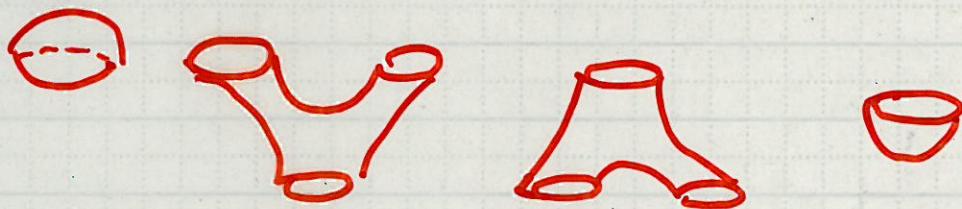
(Free on a Frobenius algebra)

- objects      formal finite sums of circles  
( $\sim$  compact closed 1-manifolds)
- maps      (isomorphism classes of) topological  
Riemann surfaces with  
appropriate  $\partial$  (convention).

## PICTURE



## Generators



BICATEGORIES  
AND THE  
PENTAGON

# BICATEGORIES

## Weak 2-categories

1-cell composition associative wrt unit

UP TO ISOMORPHISM

$$1_f \xrightarrow{\lambda} f \quad g! \xrightarrow{\rho} g \quad (hg)f \xrightarrow{\alpha} h(gf)$$

wth COHERENCE CONDITIONS

$$\begin{array}{ccc}
 ((kh)g)f & \xrightarrow{\alpha} & (kh)(gf) \\
 \alpha_f \swarrow & & \searrow \alpha \\
 (k(hg))f & & k(h(gf)) \\
 \alpha \searrow & & \nearrow k\alpha \\
 & k((hg)f) &
 \end{array}$$

$$\begin{array}{ccc}
 (g!)f & \xrightarrow{\lambda} & g(1_f) \\
 \rho_f \searrow & & \swarrow g\lambda \\
 & gf &
 \end{array}$$

(Two more are natural but Kelly showed they are redundant.)

## MAC LANE COHERENCE

- All diagrams constructed from  
 $\lambda, \rho, \alpha$   
commute. (A property of term rewriting!)  
(So  $\lambda, \rho, \alpha$  produce no essential ambiguity.)
- Any bicategory is 'equivalent' to a 2-category.  
So WHY BOTHER?

**WARNING** Not true at dimension 3 and above; nor with extra structure.

Also  $\alpha$  has meaning.

## DUAL GROUP ALGEBRA

Let  $\mathcal{F}G = \mathbb{C}^G$  be the algebra of  $\mathbb{C}$ -valued functions on the finite group  $G$  with pointwise multiplication.

(Dual to group algebra  $\mathbb{C}[G]$ .)

Sufficient to consider 1-dimensional representations of the form

$$\mathbb{C}g \text{ for } g \in G$$

with, for  $\phi \in \mathcal{F}G$

$$\phi \cdot g = \phi(g) g.$$

## TENSORING FG-MODULES

comes from comultiplication on  $FG$

(that is from multiplication on  $\mathbb{C}[G]$ ).

So we have

$$\mathbb{C}_g \otimes \mathbb{C}_h \cong \mathbb{C}_{gh}.$$

And we might as well take equality.

So choosing an associativity constraint  
is choosing

$$C_{abc} = ((\mathbb{C}_a \otimes \mathbb{C}_b) \otimes \mathbb{C}_c \xrightarrow{\alpha_{abc}} \mathbb{C}_a \otimes (\mathbb{C}_b \otimes \mathbb{C}_c)) = C_{abc}$$

such that ...

# MAC LANE PENTAGON IN FG-MODULES

$$\begin{array}{ccc}
 ((a \otimes b) \otimes c) \otimes d & \xrightarrow{\alpha_{ab,c,d}} & (a \otimes b) \otimes (c \otimes d) \\
 \downarrow \alpha_{a,b,c} & & \downarrow \alpha_{a,b,cd} \\
 (a \otimes (b \otimes c)) \otimes d & & a \otimes (b \otimes (c \otimes d)) \\
 \downarrow \alpha_{a,b,c,d} & & \uparrow \alpha_{b,c,d} \\
 a \otimes ((b \otimes c) \otimes d) & &
 \end{array}$$

$$\alpha_{a,b,cd} \cdot \alpha_{ab,c,d} = \alpha_{b,c,d} - \alpha_{a,bc,d} \cdot \alpha_{a,b,c}$$

Take logs and rearrange!

## COHOMOLOGY OF GROUPS

Let  $G$  a group.

Given an abelian group  $A$  set

$$C_n = \text{Maps}(G^n, A)$$

and define

$$d : C_n \longrightarrow C_{n+1}$$

by

$$d\phi(g_0, \dots, g_n)$$

$$\begin{aligned} &= \phi(g_0, \dots, g_{n-1}) - \phi(g_0, \dots, g_{n-1}, g_n) + \phi(g_0, \dots, g_{n-2}, g_{n-1}, g_n) \\ &\quad + \dots + (-1)^n \phi(g_0, g_1, \dots, g_n) + (-1)^{n+1} \phi(g_1, \dots, g_n) \end{aligned}$$

## FIRST APPROXIMATION

An associativity constraint is a  
3 co-angle i.e. a  
member of  $Z_3(G; \mathbb{C}^*)$ .

But do the boundaries mean  
anything?

$B_3(G; \mathbb{C}^*)$  gives monoidal isomorphism  
with the trivial identity constraint.

## SECOND APPROXIMATION

An up-to-isomorphism associativity  
constraint is an element of  
 $H_3(G, \mathbb{C}^*)$ .

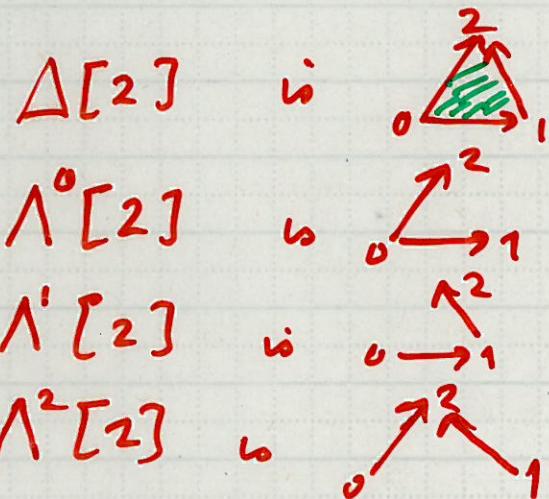
HIGHER  
DIMENSIONAL  
CATEGORY  
THEORY

## HORNS

The  $n$ -simplex  $\Delta[n]$  has  $n+1$  vertices and opposite each vertex  $k$  with  $0 \leq k \leq n$  there is an  $n-1$  dimensional face.

The  $k^{\text{th}}$  horn  $\Lambda^k[n]$  is the union of the  $(n-1)$ -dimensional faces other than the  $k^{\text{th}}$ .

Example



## KAN COMPLEXES

A simplicial set is a Kan complex if and only if every horn has a filler

$$\begin{array}{ccc} \Lambda^k[n] & \longrightarrow & X \\ \downarrow & \ddots \nearrow \dashv & \\ \Delta[n] & & \end{array}$$

(The fundamental category of a Kan (i.e. via left adjoint) complex is already a groupoid.)

## QUASI-CATEGORIES

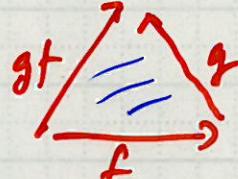
$\Delta^k[n]$  is an inner horn when  $0 < k < n$ .

A simplicial set  $X$  is a quasi-category if and only if all inner horns have fillers.

$$\text{inner } \Delta^k[n] \longrightarrow X$$
$$\Delta^n[n]$$

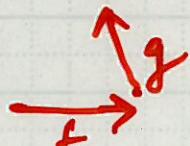
(In Boardman + Vogt : Joyal has systematically developed category theory for quasi categories.)

## ONE DIMENSIONAL MEANING

$\Delta[2]$  in  $X \in [\Delta^0, \text{Set}]$  is 

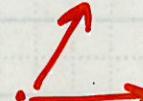
where the one-cell is a 'reason' why

$$g \cdot f \cong gf.$$

$\Lambda'[2]$  in  $X$  is 

is to find a composite and reason why.

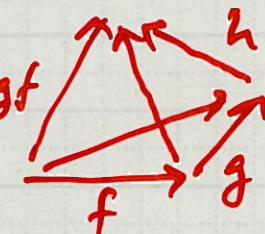
There is no need to try to fill

$\Lambda^0[2]$  in  $X$  

$\Lambda^2[2]$  in  $X$  .

## TWO DIMENSIONAL MEANING

$\Delta[3]$  in  $X$  is



which is

reasons why

$$g \cdot f \simeq gf$$

$$h \cdot g \simeq hg$$

$$h \cdot gf \simeq hgf$$

$$hg \cdot f \simeq hgf$$

plus a reason why  
the different ways to  
compose are  
equivalent.

$\Lambda'[3]$  in  $X$  misses  $h \cdot gf \simeq hgf$  and  
one expects to compose to get that.

$\Lambda^{\circ}[3]$  in  $X$  misses  $h \cdot g \simeq hg$  and  
there's no evident way to reconstruct it.

AND SO ON !?

# WEAK HIGHER CATEGORIES

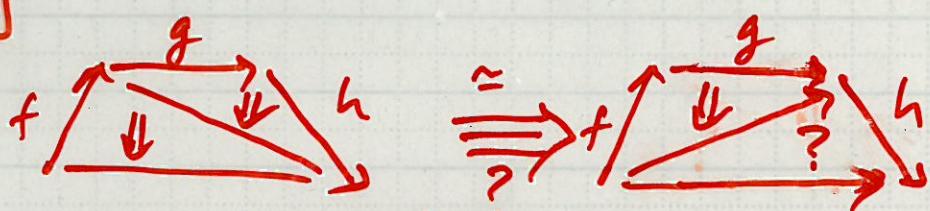
(style of Combinatorial Sets (Roberts: Street, Verity))

Quasi-categories are weak higher categories with all higher cells  $\geq 2$  equivalences.

IDEA

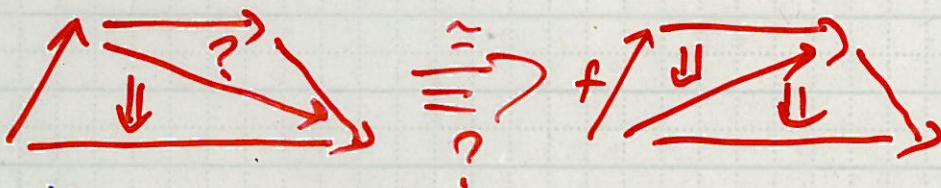
(party discipline.)

$\wedge^1[3]$



This ok if  $f \uparrow \stackrel{\simeq}{\Rightarrow} g$  is an equivalence.

$\wedge^0[3]$



This ok if  $\Delta \uparrow \stackrel{\simeq}{\Rightarrow}$  is an equivalence

and

$f \uparrow \stackrel{\simeq}{\Rightarrow}$  is an equivalence

## WEAK ALGEBRA

Operations given up to equivalences  
which are themselves up to  
equivalence and so on.

(Grothendieck's idea!)

But under what circumstances  
should you be able to "fill in"  
i.e. solve equations? And what  
is the combinatorial information  
implicit?